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## The abc conjecture

Jeffrey Paul Wheeler University of Tennessee

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To the Graduate Council:

I am submitting herewith a thesis written by Jeffrey Paul Wheeler entitled "The abc conjecture." I have examined the final electronic copy of this thesis for form and content and recommend that it be accepted in partial fulfillment of the requirements for the degree of Master of Science, with a major in Mathematics.

Pavlos Tzermias, Major Professor

We have read this thesis and recommend its acceptance:

Accepted for the Council: Carolyn R. Hodges

Vice Provost and Dean of the Graduate School

(Original signatures are on file with official student records.)

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*6.<&, ML 7* 

Acceptance for the Council:

Vice Provost and Dean of **Graduate Studies** 

## **The abc Conjecture**

A Thesis Presented for the Master of Science Degree The University of Tennessee, Knoxville

 $\mathbb{R}^d$  .

Jeffrey Paul Wheeler August 2002

# Thesis **�00�**  *. · 4485*

## **Dedication**

*This thesis is dedicated to* 

## *My Parents*

#### *whose loving sacrifices*

*ensured me the best possible education.* 

'

**" ... for we were born only yesterday and know nothing, and our days on earth are but a shadow."** 

**Job 8:9 (NIV}** 

### **Acknowledgements**

A great number of individuals are owed a debt of gratitude for their contributions to any academic success I have attained. The following are but a few of those individuals.

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My family deserves recognition for their continual support - my mother, Ruth, father, Jeffrey, and brother, David, as well as other extended family memebers. As well I wish to thank my in-laws, Jack and Cheryl Stankus, mostly for their daughter, but as well for their assistance. In particular, my wife, Jamie, is to be praised for her love and soft encouragement. I am blessed to have her in my life.

### **Abstract**

**The problem of the abc Conjecture is stated and various consequences are established. Other known consequences are stated without proof. Topics supporting belief that the abc Conjecture is true are discussed. The idea of good abc triples is defined and all known good triples are stated. Some computational computer work verifying these values is discussed. This is the first time that such brute force computations have been published.** 

## **Contents**



# **List of Tables**

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# **Chapter 1 Introduction**

**On its own the abc Conjecture merits much admiration. As is often the case with some of the more intriguing problems of Number Theory, the abc Conjecture is easy to state but yet difficult to verify. Unlike most other Number Theory problems, though, this conjecture has many fascinating applications; one of which is a version of one of the subject's most celebrated problems.** 

**Pierre de Fermat {1601 - 1665) stated his "Last Theorem" 1 in the margin of his copy of Diophantus's** *Arithmetica* **in 1637. In one of the boldest claims by one of the brightest individuals in the history of mathematics, Fermat wrote that he had a proof, but that he did not have enough room to write it in the margin. It is very likely that his proof was incomplete. Nonetheless, his innocent enough statement incited hundreds of capable (and not so capable) individuals into feverish work for over three and one-half centuries. These individuals made great accomplishments in mathematics; the development of Modern Algebra being one of the foremost. This intriguing chapter in mathematics' history came to a close in 1993 with the work of Andrew Wiles. The significance of this is best summarized in a comment by John Fraleigh regarding Wiles' proof of Fermat's Last Theorem:** *"One wonders, with the pace of science today, whether any mathematician could now make a mathematical conjecture whose status (true, false, or undecidable) could not be established, despite intense effort by the best mathematicians, for another* **350** *years".* 

**Though it may be the case that the abc Conjecture is one such conjecture, it is too early to tell. As we will see, much has been accomplished, yet the conjecture's certainty or uncertainty is not in sight. The interesting connection, though, is that the abc Conjecture implies a weaker (yet significant) form of Fermat's Last Theorem (see Conjecture 3.1).** 

**The abc Conjecture was posed in 1985 by both J. Oesterle and D.W. Masser. Oesterle was motivated by a conjecture of Szpiro regarding elliptic curves. A little later, Masser was motivated by considering an analogous statement over Z of Mason's Theorem for polynomials. We will see both Szpiro's conjecture (Conjecture 3.4) and Mason's Theorem (Theorem 2.1).** 

 $1x^n + y^n = z^n$  has no nontrivial solutions in Z for  $n \geq 3$ .

## Chapter 2

## **The Problem Stated**

#### **2.1 The abc Conjecture**

**First we begin with a defintion:** 

**Definition 2.1 (The radical of a positive integer).** 

*For*  $n \in \mathbb{P}$ , *suppose*  $n = p_1^{e_1} \cdots p_k^{e_k}$  where the  $p_i$ 's are distinct prime numbers and the  $e_i$ 's *are positive integers. We then define the radical of n to be:* 

$$
r(n) = p_1 \cdots p_k \text{ with } r(1) = 1.
$$

In other words,  $r(n)$  is the greatest square-free factor of n.

**Now we concern ourselves with the hypotheses. We will be considering non-trivial triples of integers**  $(a, b, c)$  such that  $a + b = c$  and  $gcd(a, b, c) = 1$ . Obviously, any sum of the form  $a + b = c$  can be rearranged so that  $a, b, c > 0$ , hence we will assume all elements of our **triples are positive.** 

**Oesterle originally stated the conjecture in the form** 

$$
L = L(a, b, c) = \frac{\log \max(|a|, |b|, |c|)}{\log \tau(abc)} = \frac{\log c}{\log \tau(abc)}
$$

**and considered whether the** *L's* **are bounded. We will consider this more extensively in Chapter 5.1.** 

Masser refined the statement into its more common form, namely: for each  $\varepsilon > 0$  there exists a positive universal constant  $\mu(\varepsilon)$ <sup>1</sup> such that

$$
\max(|a|,|b|,|c|) = c \leq \mu(\varepsilon) r(abc)^{1+\varepsilon}.
$$

**We now state two Lemmas that will be repeatedly quite useful.** 

**Lemma 2.1.**  Under the hypotheses of the abc Conjecture,  $r$  is a muliplicative function  $^2$ .

<sup>&</sup>lt;sup>1</sup>The literature commonly refers to this constant as  $C(\varepsilon)$ .

<sup>&</sup>lt;sup>2</sup>By definition of a multiplicative function, it is already the case that the elements involved have gcd = **1. The redundancy is for emphasis.** 

*Proof.*  **Obvious.** 

**Lemma 2.2.**  *For all*  $n \in \mathbb{P}$ ,  $r(n) \leq n$ .

*Proof.*  **Obvious** 

It is worthwhile to emphasize the importance of the  $\varepsilon$  in Masser's version of the abc Con**jecture. We will do this by using an example developed by Wojtek Jastrzebowski and Dan**  Spielman as reported by Serge Lang  $[8]$ . We show that there does not exist a  $\mu$  such that  $c \leq \mu \cdot r(abc)$  for all  $a, b$ , and  $c$  meeting the hypotheses.

For an example, consider  $a_n = 3^{2^n} - 1$ ,  $b_n = 1$ , and  $c_n = 3^{2^n}$  where  $n \in \mathbb{P}$ . Note that the **values meet the conditions of the hypotheses of the abc Conjecture. First,** 

**Claim 2.1.**   $2^n|(3^{2^n}-1)$ 

*Proof.*  For  $n = 1$ ,  $2|(3^2 - 1)$ . Assume true for  $k$ , i.e.  $2^k|(3^{2^k}-1)$ . So

$$
3^{2^{k+1}} - 1 = 3^{2^k \cdot 2} - 1 \tag{2.1}
$$

$$
= (3^{2^k})^2 - 1 \tag{2.2}
$$

$$
= (3^{2^k} - 1)(3^{2^k} + 1) \qquad \qquad \text{(difference of two squares)} \tag{2.3}
$$

**Since** 

$$
2^{k} \left| \left( 3^{2^{k}} - 1 \right) \right| \qquad \qquad \text{(induction hypothesis)} \tag{2.4}
$$

**and** 

$$
2|(3^{2^k}+1)
$$
 (viz.,  $3^{2^k}+1$  is even) (2.5)

**then** 

$$
2^{k+1}|(3^{2^{k+1}}-1). \t\t(2.6)
$$

**Hence, by induction, the claim is established.** 

#### **Proposition 2.1.**

*The c in the abc Conjecture is essential.* 

**□** 

**□** 

*Proof.* 

For contradiction, assume there exists  $\mu$  such that  $c_n \leq \mu \cdot r(a_n b_n c_n)$  for the above conditions.

So

$$
\max(|a_n|, |b_n|, |c_n|) = 3^{2^n} \tag{2.7}
$$

$$
\leq \mu \cdot r(a_n b_n c_n) \tag{2.8}
$$

$$
= \mu \cdot r([3^{2^n} - 1] \cdot 1 \cdot 3^{2^n})
$$
\n(2.9)

$$
= \mu \cdot 3 \cdot r(3^{2^{n}} - 1) \qquad \qquad \text{(by Lemma 2.1)} \tag{2.10}
$$

$$
= \mu \cdot 3 \cdot r \left(2^n \cdot \frac{3^2 - 1}{2^n}\right) \qquad \text{(by Claim 2.1)} \tag{2.11}
$$

$$
\leq \mu \cdot 3 \cdot 2 \cdot \left(\frac{3^{2^n} - 1}{2^n}\right) \quad \text{(by Lemma 2.2)} \tag{2.12}
$$

(One may regard the fraction in statement (2.12) as the product of all factors of  $3^{2^n} - 1$ different from 2).

Multiplying both sides by  $2^n$  and dividing by  $3^{2^n}$ :

$$
2^n \le \mu \cdot 6 \cdot \frac{3^{2^n} - 1}{3^{2^n}}.
$$
\n(2.13)

Letting  $n \to \infty$  the inequality fails and we get the contradiction establishing the necessity of the  $\varepsilon$ .

**□** 

Before we close this section, we state a remark that the reader may find useful while contemplating later material.

**Remark 2.1.**  *In the abc Conjecture,*  $\mu(\varepsilon)$  *varies inversely with the choice of*  $\varepsilon$ *.* 

#### **2.2 The Polynomial Analogue of the abc Conjecture**

Before we consider some of the consequences of the abc Conjecture, let us take a look at one of the conjecture's influences. Recall from the Introduction that Mason's Theorem inspired Masser. Hence we shall state Mason's Theorem.

First, a definition:

#### **Definition 2.2 {The radical of a polynomial).**

Let  $p(t)$  be a polynomial whose coefficients belong to an algebraically closed field of charac*teristic* O. *Put* 

 $n_0(p) =$  *the number of distinct zeros of p(t).* 

*In other words,*  $n_0(p)$  *counts the zeros of*  $p(t)$  *by giving them each a multiplicity of one.* 

**With this definition, we may now state:** 

#### **Theorem 2.1 (Mason's Theorem). <sup>3</sup>**

*Let a(t), b(t), and c(t) be polynomials whose coefficients belong to an algebraically closed field of characteristic 0. Suppose a(t), b(t), and c(t) are relatively prime and that*  $a(t) + b(t) =$ *c(t). Then* 

$$
\max \deg\{a(t), b(t), c(t)\} \leq n_0(a(t) \cdot b(t) \cdot c(t)) - 1.
$$

*Proof.*  **We have** 

$$
a + b = c \tag{2.14}
$$

**Dividing both sides by c:** 

$$
\frac{a}{c} + \frac{b}{c} = 1\tag{2.15}
$$

Putting  $f = \frac{a}{c}$  and  $g = \frac{b}{c}$ , we have:

$$
f + g = 1 \tag{2.16}
$$

**Differentiating we get:** 

$$
f' + g' = 0 \tag{2.17}
$$

**Rewrite as:** 

$$
\frac{f'}{f} \cdot f + \frac{g'}{g} \cdot g = 0 \tag{2.18}
$$

Ğ.

$$
\frac{g'}{g} \cdot g = -\frac{f'}{f} \cdot f \tag{2.19}
$$

So

$$
\frac{g}{f} = \frac{-\frac{f'}{f}}{\frac{g'}{g}}\tag{2.20}
$$

Observe that  $a = f \cdot c$  and  $b = g \cdot c$ , hence

$$
\frac{b}{a} = \frac{g}{f} \tag{2.21}
$$

**<sup>3</sup> It is essential that the reader realizes the similarities between Mason's Theorem and Masser's version of the abc Conjecture.** 

Substituting (2.20) into (2.21)

$$
\frac{b}{a} = \frac{-\frac{f'}{f}}{\frac{g'}{g}}\tag{2.22}
$$

Now suppose  $R(t)$  is a rational function with  $\rho_i$  representing the distinct roots of the numerator and denominator. Then

$$
R(t) = \prod_{i} (t - \rho_i)^{q_i} \text{ where the } q_i \in \mathbb{Z}
$$
 (2.23)

 $\begin{cases} > 0 & \text{if } t - \rho_i \text{ is in the numerator} \\ < 0 & \text{if } t - \rho_i \text{ is in the denominator} \end{cases}$ Notice:  $q_i$  is the multiplicity of the root  $\rho_i$  where  $q_i \n\begin{cases} 0 & \text{if } t - \rho_i \text{ is in the denominator} \end{cases}$ (2.24)

Thus

$$
R'(t) = \sum_{i} q_i \cdot \frac{R(t)}{t - \rho_i} \tag{2.25}
$$

**Hence** 

$$
\frac{R'(t)}{R(t)} = \sum_{i} \frac{q_i}{t - \rho_i} \tag{2.26}
$$

The advantage of (2.26) is that the multiplicity of each distinct root is now exactly one. Now suppose

$$
a(t) = \prod_i (t - \alpha_i)^{m_i}, \ b(t) = \prod_j (t - \beta_j)^{n_j}, \text{ and } c(t) = \prod_k (t - \gamma_k)^{r_k} \qquad (2.27)
$$

Then by (2.22) and (2.26),

$$
\frac{b}{a} = -\frac{\frac{f'}{f}}{\frac{g'}{g}} = -\frac{\sum_{i} \frac{m_i}{t - \alpha_i} - \sum_{k} \frac{r_k}{t - \gamma_k}}{\sum_{j} \frac{n_j}{t - \beta_j} - \sum_{k} \frac{r_k}{t - \gamma_k}}
$$
(2.28)

A common denominator for the numerator and denominator of (2.28) is (since *a, b,* and c are relatively prime)

$$
D(t) := \prod_{i}(t - \alpha_i) \cdot \prod_{j}(t - \beta_j) \cdot \prod_{k}(t - \gamma_k)
$$
\n(2.29)

where

$$
\deg(D(t)) = n_0(abc) \tag{2.30}
$$

Now we make the observation that

$$
\deg\left(\frac{f'}{f}\right) = -\infty, \text{ if } b \equiv 0; \deg\left(\frac{g'}{g}\right) = -\infty, \text{ if } a \equiv 0; \tag{2.31}
$$

**and** 

$$
\deg\left(\frac{f'}{f}\right) = \deg\left(\frac{g'}{g}\right) = -1 \text{ if } a \text{ nor } b \equiv 0. \tag{2.32}
$$

 $\tilde{\mathcal{O}}$ 

**Hence** 

$$
\deg\left(D \cdot \frac{f'}{f}\right) \text{ and } \deg\left(D \cdot \frac{g'}{g}\right) \le n_0(abc) - 1. \tag{2.33}
$$

**(Note that in (2.31), (2.32) , and (2.33) we needed the hypothesis that the polynomials have coefficients in a field of characteristic 0. This will also be used in (2.37) , (2.38) , and (2.39).)** 

**By (2.22) we get** 

$$
\frac{b}{a} = \frac{-D \cdot \frac{f'}{f}}{D \cdot \frac{g'}{g}}\tag{2.34}
$$

**Hence** 

$$
-a \cdot \left(D \cdot \frac{f'}{f}\right) = b \cdot \left(D \cdot \frac{g'}{g}\right) \tag{2.35}
$$

 $Since (a,b) = 1$ 

$$
a \mid \left( D \cdot \frac{g'}{g} \right) \tag{2.36}
$$

**Thus by (2.33)** 

$$
\deg(a) \le n_0(abc) - 1 \tag{2.37}
$$

**A similar argument yields** 

$$
\deg(b) \le n_0(abc) - 1 \tag{2.38}
$$

**As well** 

$$
\deg(c) \le \max\left\{\deg(a), \deg(b)\right\} \tag{2.39}
$$

**So, by (2.37) , (2.38) , and (2.39) :** 

$$
\max \deg\{a(t), b(t), c(t)\} \le n_0(a(t) \cdot b(t) \cdot c(t)) - 1. \tag{2.40}
$$

**□** 

**Having established Mason's Theorem we get** 

#### **Corollary 2.1 (Fermat's theorem for polynomials).**

Let  $x(t)$ ,  $y(t)$ , and  $z(t)$  be relatively prime polynomials whose coefficients belong to an al*gebraically closed field of characteristic* **O** *such that at least one of them has degree* > **0.**  *Then* 

$$
x(t)^n + y(t)^n = z(t)^n
$$

*has no solution for*  $n \geq 3$ *.* 

*Proof.* 

**By Mason's Theorem we have** 

$$
\deg(x(t)^n) = n \cdot \deg(x(t)) \leq \deg(x(t)) + \deg(y(t)) + \deg(z(t)) - 1.
$$

By successively replacing the  $x(t)$  on the LHS with  $y(t)$  and  $z(t)$  and summing we get

$$
n[\deg(x(t) + \deg(y(t)) + \deg(z(t))] \leq 3[\deg(x(t)) + \deg(y(t)) + \deg(z(t))] - 3
$$

This is an obvious contradiction for  $n \geq 3$ .

*Remark 2.2.* 

*Fermat's theorem for polynomials fails if char*  $p > 0$ *.* 

For an example, let  $f(x) = x + 1$ ,  $g(x) = x$ , and  $h(x) = 1$  be polynomials whose coefficients are in a field of char  $p > 0$ . Then  $f(x)^p = g(x)^p + h(x)^p$ .

**□** 

## Chapter 3

## **Consequences of the abc Conjecture**

#### **3.1 Specific Consequences**

**This chapter states some of the conjecture's fascinating implications. We begin with one of the more interesting ones. As stated in the Introduction, the abc Conjecture implies a**  weaker form of Fermat's Last Theorem. This is due to the  $\mu(\varepsilon)$ .

**Conjecture 3.1 (The Asymptotic Fermat Problem). <sup>1</sup>** *Then there exists*  $N \in \mathbb{Z}$  *such that for*  $n > N$ ,

 $x^n + y^n = z^n$ ,

*where*  $gcd(x, y, z) = 1$ *, has only trivial solutions in the integers.* 

**Theorem 3.1.**  *The abc Conjecture implies the Asymptotic Fermat Problem.* 

Notation  $(x(t) \ll y(t))$ . *We will use the symbol*  $\ll$  *to mean the following: For functions x(t) and y(t)* 

 $x(t) \ll y(t)$  means  $\exists C \in \mathbb{R}$ ,  $C > 0$  such that  $x(t) \leq C \cdot y(t)$  for all t.

*Another way to state this is that in big oh notation*  $x(t) \ll y(t)$  means  $x(t) = O(y(t))$ .

*Proof of Theorem 3.1.* 

**Again, we may make the appropriate rearrangements in the sum so that all integers are positive.** 

By **the abc Conjecture:** 

$$
|x^n| \le \mu\left(\frac{\varepsilon}{3}\right) \cdot r(xyz)^{1+\frac{\varepsilon}{3}} \ll |xyz|^{1+\frac{\varepsilon}{3}} \tag{3.1}
$$

$$
|y^n| \leq \mu\left(\frac{\varepsilon}{3}\right) \cdot r(xyz)^{1+\frac{\varepsilon}{3}} \ll |xyz|^{1+\frac{\varepsilon}{3}} \tag{3.2}
$$

**<sup>1</sup> In the wake of Andrew Wiles' accomplishment, this conjecture may be labeled a corollary or, more aptly, an academic corollary.** 

**and** 

$$
|z^n| \le \mu\left(\frac{\varepsilon}{3}\right) \cdot r(xyz)^{1+\frac{\varepsilon}{3}} \ll |xyz|^{1+\frac{\varepsilon}{3}}.
$$
\n(3.3)

**Hence** 

$$
|x^n|\cdot|y^n|\cdot|z^n|=|xyz|^n\ll(|xyz|^{1+\frac{\epsilon}{3}})^3=|xyz|^{3+\epsilon}.\tag{3.4}
$$

Thus for  $|xyz| > 1$  we get *n* bounded. Otherwise,  $|xyz| \le 1$  and at least one of the integers **must be 0.** □

It is worthwhile to note the role of the  $\mu(\epsilon)$  in the previous proof. In particular, our choice of  $\varepsilon$  determines the value of the  $N$ .

**The abc Conjecture also implies the following classical conjecture. Before we state the conjecture, we establish the necessary definition.** 

#### **Definition 3.1 (Wieferich Condition).**

*A* prime  $p \in \mathbb{Z}$  satisfies the **Wieferich Condition** iff  $2^{p-1} \not\equiv 1 \mod p^2$ .

**Conjecture 3.2 (Infinity of Primes Satisfying the Wieferich Condition).**  There exist infinitely many primes p satisfying the Wieferich Condition.

#### **Theorem 3.2.**

*The abc Conjecture implies that an infinity of primes satisfy the Wieferich Condition.* 

It will be helpful to employ the set  $S := \{p \mid p \text{ is prime and } 2^{p-1} \neq 1 \text{ mod } p^2\}$ . But before **we prove the theorem, we first establish the following claim:** 

**Claim 3. 1.**  *Let*  $n \in \mathbb{P}$  and  $p$  be a prime such that  $2^n \equiv 1 \mod p$  but  $2^n \not\equiv 1 \mod p^2$ . Then  $p \in S$ .

*Proof of Claim 3.1.*  Put  $d = \text{ord}(2)$  in  $U(\mathbb{Z}/p\mathbb{Z})$  where  $|U(\mathbb{Z}/p\mathbb{Z})| = p - 1$  Hence  $d | (p - 1)$  and  $d | n$ . *Write*  $n = dr$  for some  $r \in \mathbb{Z}$ . So  $2^n \not\equiv 1 \mod p^2 \Rightarrow 2^d \not\equiv 1 \mod p^2$ . *Now write*  $p - 1 = dm$  for some  $m \in \mathbb{Z}$ .  $d \leq p - 1$ ,  $\therefore m \leq p - 1 \leq p$ . Also, p prime  $\Rightarrow$   $(m, p) = 1.$  $\Rightarrow$   $(m, p) = 1.$ <br>*Since*  $d = \text{ord}(2)$  *in U(Z/pZ),*  $2^d \equiv 1 \text{ mod } p.$ Hence  $\exists k \in \mathbb{Z}$  such that  $pk = 2^d - 1 \Leftrightarrow 2^d = 1 + pk$ . *Since*  $2^d \neq 1 \mod p^2$ , *it follows that*  $p \nmid k$ . *Since*  $2^a \neq 1 \text{ mod } p^2$ , it follows that  $p \nmid k$ .<br> *So*  $2^{p-1} = (2^d)^m \equiv 1^m \text{ mod } p = 1 \text{ mod } p$ . *But* 

$$
2^{p-1} = 2^{dm} = (2^d)^m = (1 + pk)^m = \binom{m}{0} 1^m (pk)^0 + \binom{m}{1} 1^{m-1} (pk)^1 + \underbrace{\sum_{i=2}^m \binom{m}{i} 1^{m-i} (pk)^i}_{\text{divisible by } p^2}
$$

$$
\equiv 1 + mpk (mod p2)
$$
  

$$
\not\equiv 1 (mod p2)
$$
 (since  $p \nmid k$  and  $p \nmid m$ )

 $\therefore$   $p \in S$  and the claim is established.

**□** 

**The following proof is due to Silverman.** 

#### *Proof of Theorem 3.2.*

Suppose S is finite. Write  $2^n - 1 = u_n v_n$  where  $\forall p_i \mid u_n, p_i \in S$  and each  $p_k \mid v_n, p_k \notin S$ . Suppose 5 is finite. Write  $2^n - 1 = u_n v_n$  where  $\forall p_i \mid u_n, p_i \in S$  and each  $p_k \mid v_n, p_k \notin S$ .<br>
S finite  $\Rightarrow u_n$  is bounded. Suppose  $p \mid v_n$ . By the claim,  $2^n \equiv 1 \mod p^2$ , i.e.  $p^2 \mid (2^n - 1)$ .  $\therefore$   $p^2 \mid u_n v_n$  (since  $2^n - 1 = u_n v_n$ ). But  $p \nmid u_n$ ,  $\therefore$   $p^2 \mid v_n$ . Since  $(2^n - 1) + 1 = 2^n$ , by the abc Conjecture

$$
| 2n - 1 | = unvn \leq \mu(\varepsilon) \cdot r(unvn)1+\varepsilon
$$
 (3.5)

$$
\leq \mu(\varepsilon) \cdot (u_n v_n^{\frac{1}{2}})^{1+\varepsilon} \tag{3.6}
$$

$$
\ll (u_n v_n^{\frac{1}{2}})^{1+\varepsilon} \tag{3.7}
$$

$$
=u_n^{1+\varepsilon}\cdot v_n^{\frac{1+\varepsilon}{2}}\tag{3.8}
$$

$$
\ll v_n^{\frac{1+\epsilon}{2}}.\tag{3.9}
$$

**Therefore** 

$$
u_n v_n \ll v_n^{\frac{1+\varepsilon}{2}}.\tag{3.10}
$$

Multiplying both sides by  $v_n \frac{-1-\epsilon}{2} \cdot u_n^{-1}$ 

$$
v_n^{\frac{1-\epsilon}{2}} \ll u_n^{-1} \tag{3.11}
$$

**Hence** 

$$
v_n \ll u_n^{\frac{2}{\epsilon - 1}} \tag{3.12}
$$

$$
\Rightarrow \text{ a finite number of } v_n \tag{3.13}
$$

contradiction as 
$$
n \to \infty
$$
. (3.14)

□

**Regarding the Wieferich Condition, there are only two known exceptions. Moreover, by**  Regarding the Wieferich Condition, there are only two known exceptions. Moreover, by the Lang-Trotter conjectures, the probibilty that  $2^{p-1} \equiv 1 + pk (modp^2)$  for a fixed residue class *k* mod *p* should be  $O(\frac{1}{p})$ . Hence, for fixed x, the number of primes  $p \leq x$  such that class  $\kappa$  mod  $p$  should be  $O(\frac{\pi}{p})$ . Hence, for fixed x, the number of primes  $p \leq x$  such that  $2^{p-1} \equiv 1 + pk (modp^2)$  should be  $O(\sum_{p \leq x} \frac{1}{p}) = O(\log \log x)$ ; i.e. most primes should satisfy **the Wieferich Condition.** 

**Conjecture 3.3 {Hall's Original Conjecture).**  Let u, v be relatively prime<sup>2</sup> nonzero integers such that  $u^3 - v^2 \neq 0$ . Then

$$
|u^3-v^2| \gg |u|^{\frac{1}{2}-\varepsilon}.
$$

**Theorem 3.3.** 

*The abc Conjecture implies Hall's Original Conjecture.* 

**The following proof is due to Lang.** 

**<sup>2</sup>Originally the assumption that u and v be relatively prime was not made. This is remedied by removing any common factor and then proceeding as dictated in the proof.** 

*Proof.* 

Note that we could equivalently state that  $v^2 = u^3 + t$ ,  $t \in \mathbb{Z}$ , has a bound for  $t^3$ . In particwhere that we could equivalently state that  $\mathcal{U} = \mathcal{U}$ ,  $\mathcal{U} = \mathcal{U}$ , and a sound for  $\mathcal{U} = \mathcal{U}$  in plants when  $\mathcal{U} = \mathcal{U}$  and  $\mathcal{U} = \mathcal{U}$ . We prove a more general statement:

Fix nonzero  $a, b \in \mathbb{Z}$  and let  $m, n \in \mathbb{P}$  be such that  $mn > m + n$ . Put

$$
a \cdot u^m + b \cdot v^n = k. \tag{3.15}
$$

Fix  $\epsilon' > 0$ . By the abc Conjecture

$$
|u|^m \ll |uv \cdot r(k)|^{1+\varepsilon'}.\tag{3.16}
$$

**A similar argument yields** 

$$
|v|^n \ll |uv \cdot r(k)|^{1+\varepsilon'}.\tag{3.17}
$$

**Without loss of generality, now suppose** 

$$
|a \cdot u^m| \le |b \cdot v^n|.\tag{3.18}
$$

**Then** 

$$
|u| \ll |v|^{\frac{n}{m}}.\tag{3.19}
$$

**Subsituting into (3.17)** 

$$
|v|^n \ll |v^{1+\frac{n}{m}} \cdot r(k)|^{1+\varepsilon'} = |v|^{(1+\frac{n}{m})(1+\varepsilon')} \cdot r(k)^{1+\varepsilon'}.
$$
 (3.20)

**Hence** 

$$
|v|^{n - (\frac{m+n}{m})(1+\varepsilon')} \ll r(k)^{1+\varepsilon'}.
$$
\n(3.21)

**Thus** 

$$
|v| \ll r(k)^{\frac{m(1+\epsilon')}{mn - (m+n)(1+\epsilon')}} \ll k^{\frac{m(1+\epsilon')}{mn - (m+n)(1+\epsilon')}}.
$$
\n(3.22)

**(3.23)** 

(Note that we needed Lemma 2.2, namely:  $r(k) \leq k$ .)

**By (3. 19) ,** 

$$
|u| \ll r(k)^{\frac{m(1+\epsilon')}{mn - (m+n)(1+\epsilon')} \cdot \frac{n}{m}}.
$$
\n(3.24)

$$
=r(k)^{\frac{n(1+\epsilon')}{mn-(m+n)(1+\epsilon')}}.\t(3.25)
$$

$$
\ll k^{\frac{m(1+\epsilon')}{mn - (m+n)(1+\epsilon')}}.
$$
\n(3.26)

**<sup>3</sup>Note that the abc Conjecture, if true, improves Baker's bound for this situation.** 

**Having established the general case, we may establish the implication of Hall's Conjecture. Pick**  $\varepsilon$  such that  $\varepsilon = \frac{12\varepsilon'}{1-5\varepsilon'}$ , i.e.  $\varepsilon' = \frac{\varepsilon}{12+5\varepsilon}$ . Put  $m = 3$  and  $n = 2$ . **By (3.26),** 

$$
|u| \ll k^{\frac{2+2\varepsilon'}{1-5\varepsilon'}} = k^{2+\frac{12\varepsilon'}{1-5\varepsilon'}}.
$$
\n(3.27)

**Thus** 

$$
|u|^{\frac{1}{2}-\frac{12\epsilon'}{1-5\epsilon'}} \ll k^{(2+\frac{12\epsilon'}{1-5\epsilon'})\left(\frac{1}{2}-\frac{12\epsilon'}{1-5\epsilon'}\right)} = k^{[1-\frac{3}{2}\cdot\frac{12\epsilon'}{1-5\epsilon'} - \left(\frac{12\epsilon'}{1-5\epsilon'}\right)^2]}.
$$
\n(3.28)

Substituting for  $\varepsilon'$ 

$$
|u|^{\frac{1}{2}-\varepsilon} \ll k^{(1-\frac{3}{2}\cdot\varepsilon-\varepsilon^2)} < k. \tag{3.29}
$$

**□** 

**Recall from the Introduction that Oesterle was inspired by a conjecture of Szpiro. Hence we shall consider Szpiro's Conjecture. First, some preliminaries.** 

**Since we are considering fields of characteristic O we may assume that our elliptic curves have Weierstrass equations of the form** 

$$
E: y^2 = x^3 - ux + v \tag{3.30}
$$

where  $u, v \in \mathbb{Z}$ . Given this we identify the *disciminant of E*, namely

$$
\Delta = 16(4u^3 - 27v^2)
$$

and  $D := 4u^3 - 27v^2$  is the discrimanant of the cubic polynomial. Also, we will want to indentify the *conductor of E*, namely for prime  $p \in \mathbb{Z}$ 

$$
c(E):=\prod\nolimits_{p}p^{f_p}
$$

**where** 

 $f_p = \begin{cases} 0 & \text{if the reduction of E is non-singular} \\ 1 & \text{if the reduction of E is multiplicative} \end{cases}$  $2 + \delta_p$  if the reduction of E is additive

and  $\delta_p$  is a bounded constant independent of the curve with  $\delta_p = 0$  if  $p \geq 5$ .

**Before we continue, it is important to observe that** 

$$
r(D) \le c(E). \tag{3.31}
$$

#### **Conjecture 3.4 ( Original Szpiro Conjecture).**

*Assuming a Weierstrass equation with D the discriminant of the cubic polynomial and c(E) the conductor of the equation, then* 

$$
|D| \ll r(D)^{6+\varepsilon} \ll c(E)^{6+\varepsilon}.
$$

*(Note that Szpiro did not include the notion of r(D) in his Conjecture.)* 

**Theorem 3.4.** 

*The abc Conjecture implies the Original Szpiro Conjecture.* 

*Proof.*  Fix  $\epsilon > 0$  and put  $\epsilon'' = \frac{1}{3}\epsilon$ . We have

$$
4u^3 - 27v^2 = D.
$$

By the abc Conjecture (in particular, our proof of Hall's Conjecture)

$$
|u|^3 \ll [(r(D))^{2+\epsilon''}]^3
$$
 by (3.27) and noting that  $r(D) < D$  (3.32)

and

$$
|v|^2 \ll [(\tau(D)^{3+\epsilon''}]^2 \qquad \qquad \text{by (3.22)} \tag{3.33}
$$

Hence

$$
|D| \ll r(D)^{6+\varepsilon} \ll c(E)^{6+\varepsilon} \qquad \text{by (3.31)} \tag{3.34}
$$

$$
\qquad \qquad \Box
$$

#### **Remark 3.1.**

*The abc Conjecture is equivalent to Szpiro 's Original Conjecture.* 

For the proof of the opposite implication, see [8].

#### **3.2 Futher Consequences**

In this section we list further consequences of the abc Conjecture without proof. For further information, see [9].

#### **Definition 3.2 (Brown Pairs).**

*Pairs of integers satisfying Brocard's Problem*  $n! + 1 = m^2$  *are called Brown Pairs.* 

#### **Theorem 3.5.**

*The abc Conjecture implies that there exist only finitely many Brown Pairs.* 

The proof of this is in [11].

For the interested reader, the above problem has been generalized to the number of integer solutions of the equations  $(x!)^n + 1 = y^m$  (see [10]) and  $x! + B^2 = y^2$  for arbitrary B (see [3]).

#### **Definition 3.3 (Powerful Numbers).**

For  $n \in \mathbb{P}$ , n is said to be a **powerful number** if for every prime p dividing n,  $p^2$  divides *n.* 

*Erdos* refers to theses numbers as k-ful numbers where the *k* plays the role of the 2 in the above definition.

Conjecture 3.5 *(Erdos* -*M ollin* -*Walsh Conjecture). There do no exist three consecutive powerful integers.* 

#### **Theorem 3.6.**

*The abc Conjecture implies that the set of triples of consecutive powerful integers is finite.* 

#### **Conjecture 3.6 {Mordell's Conjecture). 4**

*Any curve of genus larger than* l *defined over a number field K has only finitely many rational points in K.* 

**The following is due to Elkies [4] .** 

#### **Theorem 3.7.**

The abc Conjecture for number fields implies the Mordell Conjecture over an arbitrary *number field.* 

In [5] it is established that the abc Conjecture with an explict constant  $\mu(\varepsilon)$  would give **explicit bounds on the heights of rational points in Mordell's Conjecture.** 

#### **Theorem 3.8 {Roth's Theorem).**

*Fix*  $\epsilon > 0$ . For every algebraic number  $\alpha$ , the diophantine inequality

$$
|\alpha-\frac{p}{q}|<\frac{1}{q^{2+\varepsilon}}
$$

*has only finitely many solutions.* 

**In 1994, E. Bombieri (1] proved that the abc Conjecture implies a stronger version of Roth's Theorem:** 

**Theorem 3.9.** 

*The abc Conjectue implies that, for the conditions of Roth's Theorem,* 

$$
|\alpha-\frac{p}{q}|>\frac{1}{q^{2+k}}
$$

for all but a finite number of fractions  $\frac{p}{q}$  in reduced form, where  $k = C(\alpha) \cdot (\log q)^{-\frac{1}{2}}$ .  $(\log \log q)^{-1}$  for some constant  $C(\alpha)$  depending only upon  $\alpha$ .

**For the above, the reader may also see [5].** 

**The following is due to Granville [6]:** 

#### **Theorem 3.10 {Squarefree Values of Polynomials).**

*The abc Conjecture implies that for a polynomial F(x) with integer coefficients, no repeated roots, and content = 1,*  $F(n)$  *is squarefree for infinitely many integers n.* 

**In closing, we mention that the abc Conjecture also gives a way of counting squarefree values of polynomials, implies that the Dirichlet £-function has no Siegel zeros, and gives bounds for the order of the Tate-Shafarevich group. Many more implications are given in**  [9].

**<sup>4</sup>This is now a theorem after the work of G. Faltings (1984).** 

## **Chapter 4**

## **Evidence for the abc Conjecture**

**In this chapter, a theorem of C.L. Stewart and Kunrui Yu establishing a weak form of the abc Conjecture is discussed.** 

#### **4. 1 Preliminaries**

**For the following, let p be a prime number and put** 

$$
q = \begin{cases} 2 & \text{if } p > 2 \\ 3 & \text{if } p = 2. \end{cases} \tag{4.1}
$$

**As well put** 

$$
\alpha_0 = \begin{cases} \zeta_4 & \text{if } p > 2 \\ \zeta_6 & \text{if } p = 2, \end{cases}
$$
 (4.2)

where  $\zeta_m$  has the usual meaning  $e^{\frac{2\pi i}{m}}$  for  $m \in \mathbb{P}$ . Put  $K = \mathbb{Q}(\alpha_0)$  and let  $D = \Omega \cap K$ , i.e. *D* is the ring of algebraic integers in *K* (*Note:* since *K* is a cyclotomic field,  $D \equiv \mathbb{Z}[\zeta_0]$ ). For  $c = x + iy \in \mathbb{C}$ ,  $|c| = \sqrt{x^2 + y^2}$ . Let  $\alpha_1, \ldots, \alpha_n \in D$  such that  $|\alpha_i| \leq A_i$  for  $1 \leq i \leq n$ where each  $A_i \geq 4$ . Put

$$
A = \max_{1 \leq i \leq n} A_i.
$$

Let  $b_1, \ldots, b_n$  be rational integers (i.e. in  $\mathbb{Z}$ ) such that  $|b_i| \leq B$  where B is a fixed integer  $\geq$  3. For  $\alpha \in K\setminus\{0\}$ , since D is a Dedekind domain the fractional ideal  $(\alpha)D$  can be written as a unique product of prime ideals in D, i.e.  $(\alpha)D = p_1^{e_{p_1}} \cdots p_g^{e_{p_g}}$ . Define  $\text{ord}_{p_i} \alpha = e_{p_i}$ . This is the ramification index of  $p_i$ . Let  $f_p$  be the residue class degree of  $p$ . Lastly, put  $\Theta = \alpha_1^{b_1} \cdots \alpha_n^{b_n} - 1.$ 

**Given the above, we now state some essential preliminaries. These are stated without proof; the curious reader may see** [12].

**Lemma 4.1.**  *If*  $[K(\alpha_0^{1/q}, \ldots, \alpha_n^{1/q}) : K] = q^{n+1}, ord_p \alpha_j = 0$  for  $j = 1, \ldots, n$ , and  $\Theta \neq 0$ , then  $\int \text{d} \sigma \, d_p \theta \leq (c_1 n)^n p^2 \cdot \log B \cdot \log \log A \cdot \log A_1 \cdot \, \cdots \, \log A_n$ 

*where* **c1** *is an effectively computable postive number.* 

#### **Lemma 4.2.**

For  $\alpha_1, \ldots, \alpha_n \in \mathbb{P}$ , if  $[\mathbb{Q}(\alpha_1^{1/2}, \ldots, \alpha_n^{1/2}) : \mathbb{Q}] = 2^n$  and  $b_1 \cdot \log \alpha_1 + \cdots + b_n \cdot \log \alpha_n \neq 0$ , *then* 

 $|b_1 \cdot \log a_1 + \cdots + b_n \cdot \log a_n| > exp(-c_2n)^n \log B(\log \log A)^2 \log A_1 \cdots \log A_n$ 

*where c<sub>2</sub> is an effectively computable positive number.* 

#### **Lemma 4.3.**

*Let*  $\alpha_1, \alpha_2, \ldots, \alpha_n$  be prime numbers with  $\alpha_1 < \alpha_2 < \cdots < \alpha_n$ . Then

$$
[\mathbb{Q}(\alpha_1^{\frac{1}{2}}, \alpha_2^{\frac{1}{2}}, \cdots, \alpha_n^{\frac{1}{2}}) : \mathbb{Q}] = 2^n.
$$

*Let*  $q = 2$  *and*  $\alpha_0 = \zeta_4$  *or*  $q = 3$  *and*  $\alpha_0 = \zeta_6$  *as well put*  $K = \mathbb{Q}(\alpha_0)$ . *Then* 

$$
[K(\alpha_0^{\frac{1}{q}}, \alpha_1^{\frac{1}{q}}, \ldots, \alpha_n^{\frac{1}{q}}):K]=q^{n+1}
$$

*except when*  $q = 2$ ,  $\alpha_0 = \zeta_4$ , and  $\alpha_1 = 2$  and in this case

$$
[K(\alpha_0^{\frac{1}{2}},(1+i)^{\frac{1}{2}},\alpha_2^{\frac{1}{2}},\ldots,\alpha_n^{\frac{1}{2}}):K]=2^{n+1}.
$$

**Lemma 4.4.** 

*Let*  $p_1 = 2$ ,  $p_2$ , be the sequence of prime numbers in increasing order. Then  $\exists$  an effectively *computable constant*  $c_3 > 0$  *such that for every positive integer r we have* 

$$
\prod_{j=1}^r \frac{p_j}{\log p_j} > (\frac{r+3}{c_3})^{r+3}.
$$

#### **4.2 The Evidence**

**Theorem 4. 1.** 

*There exists an effectively computable constant*  $k$  *such that for all*  $a, b,$  *and*  $c \in \mathbb{P}$  *with*  $(a, b, c) = 1, c > 2, and a + b = c$ 

$$
\log c < r(abc)^{\frac{2}{3} + \frac{k}{\log \log r(abc)}}
$$

**The following proof is due to Stewart and Yu.** 

#### *Proof.*

Let  $c_4, c_5, \ldots$  denote effectively computable positive constants. Without loss of generality suppose  $a \leq b$ . Since  $a + b = c$ ,  $gcd(a, b, c) = 1$ , and  $c \geq 2$ , it follows that  $a < b < c$  and that  $r(abc) \geq 6$ . Write

$$
a = p_1^{e_1} \cdots p_t^{e_t}, b = q_1^{f_1} \cdots q_u^{f_u}, and c = s_1^{g_1} \cdots s_v^{g_v},
$$

where  $p_1, \ldots, p_t, q_1, \ldots, q_u, s_1, \ldots, s_v$  are distinct primes with  $t \geq 0, u \geq 1, v \geq 1$ , and  $e, f, g \in \mathbb{P}$ . Denote the largest prime dividing *a* by  $p_a$  except when  $a = 1$  and in this **situation simply put**  $p_a = 1$ **. Similarly denote the largest primes dividing b and c by**  $p_b$  **and** *<sup>P</sup>c* **respectively. Then for any prime p** 

$$
\max\{\text{ord}_p a, \text{ ord}_p b, \text{ ord}_p c\} \le \frac{\log c}{\log 2}.\tag{4.3}
$$

**Observe that** 

$$
\log c = \sum_{p|c} (\text{ord}_p c \cdot \log p) \le (\max_{p|c} {\{\text{ord}_p c\}}) \cdot \log r(abc). \tag{4.4}
$$

Since  $(a, b) = (a, c) = (b, c) = 1$ , for each prime *p* dividing *c*,

$$
\operatorname{ord}_{p} c = \operatorname{ord}_{p} \left( \frac{c}{-b} \right) = \operatorname{ord}_{p} \left( \frac{a}{-b} - 1 \right) \le \operatorname{ord}_{p} \left( \left( \frac{a}{b} \right)^{4} - 1 \right). \tag{4.5}
$$

**We now estimate** 

$$
\mathrm{ord}_p(\left(\frac{a}{b}\right)^4 - 1) = \mathrm{ord}_p(p_1^{4e_1} \cdots p_t^{4e_t} \cdot q_1^{-4f_1} \cdots q_u^{-4f_u} - 1)
$$

for each prime *p* dividing *c*. We do this by employing Lemma 4.1.

Put  $\Theta = \left(\frac{a}{b}\right)^4 - 1$ . If  $p = 2$ , we put  $K = \mathbb{Q}(\zeta_6)$ , while if  $p > 2$  we put  $K = \mathbb{Q}(\zeta_4)$ . Define q and  $\alpha_0$  as in statements (4.1) and (4.2). Now let  $\rho$  be a prime ideal of the ring of algebraic **integers of K lying above the prime** *p.* **Thus we have** 

$$
\mathrm{ord}_p\Theta\leq\mathrm{ord}_{\boldsymbol{p}}\Theta.
$$

For *n* in Lemma 4.1, put  $n = t + u$ . As well let  $\alpha_1, \ldots, \alpha_n$  be the primes  $p_1, \ldots, p_t, q_1, \ldots, q_u$ arranged in increasing order, except in the case when  $p > 2$  and  $\alpha_1 = 2$ . In this situation, **take**  $a_1 = 1 + i$  **instead of**  $a_1 = 2$  **and note that**  $2^4 = (1 + i)^8$ **. Since p|c and**  $(a, c) = (b, c) = 1$ we have  $\text{ord}_p\alpha_i = 0$  for  $i = 1, ..., t + u$ . Clearly  $\Theta \neq 0$ . Thus, by Lemma 4.3,

$$
[K(\alpha_0^{\frac{1}{q}}, \alpha_1^{\frac{1}{q}}, \ldots, \alpha_{t+u}^{\frac{1}{q}}):K] = q^{t+u+1}.
$$

**Now put** 

 $B = \max\{8e_1, \ldots, 8e_t, 8f_1, \ldots, 8f_u\}.$ 

**So, by (4.3),** 

$$
B \leq 8 \cdot \frac{\log c}{\log 2}
$$

**Hence by Lemma 4.1** 

$$
\text{ord}_pc \leq \text{ord}_p \Theta < (c_4 \cdot (t+u))^{t+u} \cdot p^2 \cdot \log \log c \cdot \log \log r(abc) \cdot \prod_{p|ab} \log p. \tag{4.6}
$$

Similarly if  $p|b$  then, by considering  $\text{ord}_p((\frac{c}{a})^4 - 1)$ , we have

$$
\text{ord}_p b < (c_5 \cdot (t+v))^{t+v} \cdot p^2 \cdot \log \log c \cdot \log \log r(abc) \cdot \prod_{p \mid ac} \log p \tag{4.7}
$$

**·** 

and if  $p|a$  then, by considering  $\text{ord}_p((\frac{c}{b})^4 - 1)$ , we also have

$$
\operatorname{ord}_p a < (c_6 \cdot (u+v))^{u+v} \cdot p^2 \cdot \log \log c \cdot \log \log r(abc) \cdot \prod_{p|bc} \log p. \tag{4.8}
$$

**It follows from {4.4) and {4.6) that** 

$$
\frac{\log z}{\log \log c} < (c_5 \cdot (t+u))^{t+u} \cdot p_c^2 \cdot \prod_{p|ab} \log p \cdot (\log r(abc))^2. \tag{4.9}
$$

Since  $b > \frac{c}{2}$  and  $c \geq 3$ ,

$$
\log b > \log c - \log 2 > \frac{\log c}{4}.\tag{4.10}
$$

But  $(4.4)$  holds if we replace  $c$  by  $b$ . So from  $(4.7)$ 

$$
\frac{\log c}{4\log\log z} < (c_5 \cdot (t+v))^{t+v} \cdot p_b^x \cdot \prod_{p \mid ac} \log p \cdot (\log r(abc))^2. \tag{4.11}
$$

Now either  $a > \sqrt{b}$  or  $a \leq \sqrt{b}$ . Hence

$$
\begin{cases} \text{for } a > \sqrt{b}, \quad \log a \ge \frac{1}{2} \cdot \log b > \frac{\log c}{8} \\ \text{or } a \le \sqrt{b}, \quad \log(\frac{a+b}{b}) = \log(1 + \frac{a}{b}) < \log\left(1 + \frac{1}{\sqrt{b}}\right) < \frac{1}{\sqrt{b}} < \frac{\sqrt{2}}{\sqrt{c}}. \end{cases} \tag{4.12}
$$

In the former case, we use (4.4) with *c* replaced by *a* together with (4.8) to conclude that

$$
\frac{\log c}{8 \log \log c} < (c_6 \cdot (u+v))^{u+v} \cdot p_a^2 \cdot \prod_{p|bc} \log p \cdot (\log r(abc))^2. \tag{4.13}
$$

**In the latter case,** 

$$
0 < \log \frac{c}{b} = \log \left( \frac{a+b}{b} \right) = g_1 \cdot \log s_1 + \dots g_v \cdot \log s_v - f_1 \cdot \log q_1 - \dots - f_u \cdot \log q_u. \tag{4.14}
$$

By Lemma 4.3 we may use Lemma 4.2 to obtain a lower bound for  $\log \frac{c}{b}$ . Comparing this with the upper bound given by  $(4.12)$  we again obtain  $(4.13)$  with  $c_6$  replaced by  $c_7$ . Put  $\rho = u + t + v$ . From (4.9), (4.11), (4.13), we deduce that

$$
\left(\frac{\log c}{4\log\log e}\right)^3 < (c_8 \cdot \rho)^{2\rho} \cdot (p_a p_b p_c)^2 \cdot \left(\prod_{p \mid abc} \log p\right)^2 \cdot (r(abc))^6. \tag{4.15}
$$

**By Lemma 4.4,** 

$$
\left(\frac{\rho}{c_9}\right)^{\rho} < \prod_{i=1}^{p-3} \frac{p_i}{\log p_i} < 2 \cdot \prod_{p \neq p_a, p_b, p_c} \frac{p}{\log p},\tag{4.16}
$$

**with the usual convention that the empty product is 1.** 

**Thus, by(4.15),** 

$$
\left(\frac{\log c}{4\log\log c}\right)^3 < c_{10} \rho \cdot (r(abc))^2 \cdot (\log r(abc))^{12}.
$$
 (4.17)

**Again by Lemma 4.4 we have** 

$$
c_{10}^{\rho} < (r(abc))^{\frac{c_{11}}{\log\log r(abc)}},\tag{4.18}
$$

**□** 

and the result now follows from  $(4.17)$ .

**Recently the authors improved this estimate. In [13] , the better estimate** 

$$
c < exp(c_{11} \cdot (r(abc))^{\frac{1}{3}} \cdot (\log r(abc))^3
$$

where  $c_{11}$  is an effectively computatble positive constant is established. The method em**ployed to improve the estimate is a p-adic linear independence measure for logarithms of algebraic numbers. This result, due to Yu, is an ultrametric analog of an Archimedean measure due to E.M. Matveev.** 

**A second estimate is also established. In particular, if** *a, b,* **c are relatively prime positive integers such that**  $a + b = c$  and  $c > 2$ , then

$$
c < exp(p' \cdot (r(abc))^{c_{12} \cdot \frac{\log \log \log r_*(abc)}{\log \log r(abc)}},
$$

where  $c_{12}$  is an effectively computable constant,  $r_*(abc) = \max\{r(abc), 16\}$ , and  $p' =$  $min\{p_a, p_b, p_c\}.$ 

## Chapter 5

## Good Triples Associated with the **abc Conjecture**

**In this chapter, we consider the notion of good triples. We will also state the known good triples.** 

#### **5.1 Preliminaries**

**Recall Oesterle's version of the abc Conjecture, na**m**ely, under the appropriate hypotheses, he considered** 

$$
L = L(a, b, c) = \frac{\log \max(|a|, |b|, |c|)}{\log r(abc)} = \frac{\log c}{\log r(abc)}
$$

**and asked if these** *L's* **have a bound. It is this for**m **of the abc Conjecture that we will be using for the topics of this chapter.** 

#### **Theorem 5.1.**

The abc Conjecture holds iff  $\limsup\{L\} \leq 1$ .

*Proof.*   $(\Rightarrow)$ **Assu**m**e the abc Conjecture. So** 

$$
L(abc) = \frac{\log \max\{|a|, |b|, |c|\}}{\log r(abc)}\tag{5.1}
$$

$$
\leq \frac{\log[\mu(\varepsilon) \cdot r(abc)]^{1+\varepsilon}}{\log r(abc)} \text{ by the abc Conjecture} \tag{5.2}
$$

$$
= \frac{\log \mu(\varepsilon)}{\log r(a b c)} + 1 + \varepsilon. \tag{5.3}
$$

Fix  $\varepsilon > 0$ . Put  $k = \mu(\varepsilon)$ .

We want 
$$
\frac{\log k}{\log r(\text{abc})} \leq \varepsilon
$$
 for all but finitely many triples  $(a, b, c)$  (5.4)

$$
\Leftrightarrow \log r(abc) \ge \frac{\log k}{\varepsilon} \tag{5.5}
$$

$$
\Leftrightarrow r(abc) \ge M \ := \ e^{\frac{\log k}{c}}.\tag{5.6}
$$

This holds since, by the hypotheses of the abc Conjecture, there exist only finitely many  $(a, b, c)$ 's such that  $r(a, b, c) \leq M$ .

#### $(\Leftarrow)$

Suppose  $\limsup\{L\} \leq 1$ . This is true iff

$$
\limsup \left( \frac{\log c_n}{\log r(a_n b_n c_n)} \right) \le 1 \tag{5.7}
$$

$$
\Rightarrow \frac{\log c_n}{\log r(a_n b_n c_n)} \le 1 + \varepsilon \text{ for } n \text{ large.} \tag{5.8}
$$

Then for  $n > N$  for some  $N \in \mathbb{Z}$ :

$$
c_n \le r(a_n b_n c_n)^{1+\epsilon}.\tag{5.9}
$$

Choose constants  $\mu_1(\varepsilon)$ ,  $\mu_2(\varepsilon)$ , ...,  $\mu_N(\varepsilon)$  such that

$$
c_i \leq \mu_i(\varepsilon) \cdot r(abc)^{1+\varepsilon} \text{ for all } i \tag{5.10}
$$

Let

$$
\mu(\varepsilon) = \max_{1 \le i \le N} \{ \mu_i(\varepsilon) \} \tag{5.11}
$$

Thus

$$
c_n \leq \mu(\varepsilon) \cdot r(a_n b_n c_n)^{1+\varepsilon} \text{ for all } n. \tag{5.12}
$$

□

Recall from Proposition 2.1 our choices for  $a_n$ ,  $b_n$ , and  $c_n$ , namely:

$$
a_n = 3^{2^n} - 1
$$
,  $b_n = 1$ , and  $c_n = 3^{2^n}$ .

So, for these values

$$
L_n = \frac{\log 3^{2^n}}{\log r (3^{2^n} - 1 \cdot 1 \cdot 3^{2^n})}
$$
\n(5.13)

$$
= \frac{\log 3^{2^n}}{\log 3 + \log r (3^{2^n} - 1)}
$$
(5.14)

$$
\geq \frac{\log 3^{2^n}}{\log 3 + \log 2 \cdot r \left(\frac{3^{2^n} - 1}{2^n}\right)}\tag{5.15}
$$

$$
\geq \frac{2^n \log 3}{\log 3 + \log 2 + \log (3^{2^n} - 1) - \log 2^n} \qquad \text{since } r(3^{2^n}) \leq 2 \cdot \left(\frac{3^{2^n} - 1}{2^n}\right). \tag{5.16}
$$

$$
L_n \ge \frac{2^n \log 3}{\log 3 + \log (3^{2^n} - 1) - (n - 1) \cdot \log 2}.\tag{5.17}
$$

Thus for  $n = 3$ :

$$
L_3 \ge \frac{8 \cdot \log 3}{\log 3 + \log (3^8 - 1) - 2 \cdot \log 2}
$$
 (5.18)  
\n
$$
\approx 1.255203.... \tag{5.19}
$$

In particular, 
$$
L_3 > 1
$$
.

*It is easy to see that the fraction on the RHS of inequality (5.17) increases as n gets large. Hence there are infinitely many triples*  $(a_n, b_n, c_n)$  *such that*  $L_n > 1$ *. We have just shown* 

**Theorem 5.2.**  *The abc Conjecture holds iff*  $\limsup\{L_n\} = 1$ *.* 

#### **5.2 Good Triples**

#### **Definition 5.1 (Good Triple).**

*For the abc Conjecture, we say that a triple*  $(a, b, c)$  *is a good triple if*  $L > 1.4$ .

*So, by Theorem 5.2 we get* 

#### **Corollary 5.1.**

*If the abc Conjecture holds, there are only finitely many good triples.* 

The following is the list (Table 5.1) of known good abc triples as of January 2, 2002:





*So* 





#### **Discovers of the Known Good abc Triples**



### **5.3 Computations Regarding Good Triples**

**It seems that the values of the good triples displayed in Table 5.1 were discovered by means of various algorithms. Hence a brute force approach was taken to confirm that the stated values were indeed all possible good triples over a particular interval. Initially a program was written in Matlab but was found to be too inefficient. With the aid of Joel Mejeur (now with the Department of Defense) and Michael Saum (University of Tennessee, Knoxville)., a program was written in C and then run in parallel ( using MPI) on a cluster of between 24 and 30 Intel 450 MHz Pentium III computers. Initially the program checked for good**  triples over the intervals  $1 \le a \le 100,000$  and  $a \le b \le 100,000$ . Running time for this case **was approximately four and one-half days. Note that runs covering even larger intervals are underway and results will be summarized in a future paper.** 

**Results (Good Triples for**  $1 \le a, b \le 100,000$  **(Initial Run)).** 



*(No. ·) refers to the number in Table 5.1.* 

**It is worthwhile to point out that (5.22) is the good abc triple with the smallest** *c* **value.** 

**Results (Further Good Triples for**  $1 \le a, b \le 1,000,000$  (Further Run)). <sup>1</sup>



#### **5.4 Program Listings**

```
5.4. 1 abc-mpi . c
```

```
1 ,--------------------, 
 2 /* abc-mpi.c •/<br>
3 /*<br>
4 /* Written by J. Mejeur (May 2002) •/<br>
5 /* Revised by M. Saum (July 2002) •/<br>
6 /=
8 
9 
10 
11 
12 
13 
14 
      $indude $stdlib.htinclud• <atdio.h> 
tinclud• <aath.h> 
      tinclud• <api .h> 
       tinclud• <gatopt .h> 
tinclud• <aya/t1-.h> 
 1
```
<sup>&</sup>lt;sup>1</sup>This is a work in process. With some improvements we have hope of extending the ranges to 10,000,000. **The interested reader may also see [7] regarding similar unpublished work by Joe Kanapka.** 

```
15 ■truct time•al t_atart , t_■top; 
          16 ■truct tiaezon• tz_d_y; 
17 
          18 int get _priae■(int n, un■igned lnt •pria-■); 
 19 inline unsigned int get_r(unsigned int n, unsigned int *primes, int num_primes);<br>20 unsigned int gcd(unsigned int a, unsigned int b);<br>21 double diff_time(struct timeval *, struct timeval *);
22 
 23 int aain(int argc, char uargv) 
24 un■ ign•d int a, b, c, i; 
25 double •r•; 
 26 -■ ign•d int nua_priae■; 
27 -■ 1gned int IUIX•lOOOO, ain; 
28 -■ 1gnec1 int •pria-■; 
 29 double L; 
30 char tU.-[25e] ; 
31 l"ILE •fp•IIULL; 
32 
 33 int ccnmt-0; 
34 int rank-0, •iz•: 
35 
36 I• Set up MPI c-unication •/ 
                 37 MPI_C-. vorld; 
38 
 39 MPI_Iai t (large• largv) ; 
40 vorld • MPI_COIOI_WOILD; 
41 MPI_Comm_rank(vorld, &rank);<br>42 MPI_Comm_size(vorld, &size);
                 42 MPI_c ... _■ize(vorld, bize); 
43 
                 44 I• Get co..and line option■ •/ 
 45 ain-0; 
46 vbile(l) { 
 47 c•getopt(argc, argv, "l:■: ") ; 
48 it(c--1) 
 49 breu; 
50 ••itch(c) { 
50 switch(c)<br>
51 case '1':<br>
52 min=(<br>
53 break
 52 aiJa•(un■igned int) atoi(optarg) ; 
53 break; 
54 ca■• '■' : 
55 IMIX<sup>a</sup> Example int atoi(optarg);<br>
56 break;<br>
57 }
                              56 break; 
                 57 } 
 58 } 
59 
 60 /* Prime numbers calculated up to 2*mar + 1 */<br>61 /* to ensure that prime factorization can */<br>62 /* occur with c (*a+b). ***
 63 /º Primes calculated only on master (rank=0) o/<br>64 /º and sent to all slave processors via MPI. o/<br>65 if(rank==0) {<br>66 gettimeofday(&t_start,&tz_dummy);
                        67 priae■•calloc(2-..Z+l , ■izeof(un■igned int)); 
68 
  69 printt("Generating li■t of priae■\n"); 
70 ftlueb( ■tdout); 
  71 __ priae■•get _pria••C2•■az+1, priae■); 
72 printf("Jcnmd Id priae■ • \n", nua_priae■) ; 
73 ttlu■b(■tdout); 
74<br>75<br>76<br>7775 for(a•l;a<■ize;a++) { 
  76 MPI_Sand(bua_pri■••• 1, MPI_IIIT, a, 100, vorld) ; 
77 MPI_Send(priae■, nua_priae■, MPI_UNSIGIIED, a, 101 , world) ; 
  78 } 
79 
  80 ) else (<br>81 : MPI_Recv(knum_primes, 1, MPI_DNT, 0, 100, world, MPI_STATUS_IGNORE);<br>82 : primes=calloc(num_primes, sizeof(unsigned int));<br>84 } : MPI_Recv(primes, num_primes, MPI_UNSIGNED, 0, 101, world, MPI_STATUS_IGNORE
 84<br>85<br>86<br>86<br>87
                  86I• r■ array contaiA• log of radical tor ••ch nuaber 1 •. 2•■u+1 •I 
  37 / * rs array calculated on each processor. No need to send via \approx / * NPI, as send traffic could be very large. A good assumption \approx / / * is that all processors participating in the MPI VM are of \approx / 90 / * sen
  91 r■ • calloc(2-■ax+1 , ■izeof (double) ); 
92 for(i•1; 1<•2.ax+1; i++) 
93 r■Ci-1] • log(get_r(i ,pri■-■ ,nua_pri■e■ )); 
 91<br>92<br>93<br>94<br>95<br>96
                  95 it(rank••O) { 
  96 sprintf(filename, "abc.out.%d", rank);<br>97 sp=fopen(filename, "w");<br>98 getimeefdsy(&t_stop.&tz_dummy);<br>99 fprintf(fp."Time to gen primes and send = Xg\n", diff_time(&t_start,&t_stop));
100 fflush(fp);<br>101 }
101 
102 
103 / Main loop. Each processor starts with different a, increments a/<br>104 / by number of processors each time.<br>105 for(a=rank+1;a<=max;a+=size) {
106 it(rank•-0) { 
                                107 fpriAtf (fp , "rank(ld) : working on a•ld\n" ,rank, a) ;
```

```
108 ffluah(fp) ; 
 109 } 
110 
111 count++;<br>
112 if(ac=ni,<br>
113 for(
                    112 if(a<-.i.D) 
113 for(b-min;b<max;b++) {<br>
114 /* No need to calcu<br>
115 if(!(((aX2) kk
 114 /• No need to calculate gcd if both a,b divisible by same small prime */<br>115 if( ! (aX2) && !(bX2) ) || ( !(aX3) && !(bX3) ) || ( !(aX5) && !(bX5) ) || ( !(aX7) && !(bX7) ) ) } {
 1 16 if(gcd(a, b) •• 1) { 
1 17 c•a+b; 
117<br>118<br>119
                                          119 L•log ( (doubl•) c) / Cr■ [a-1] +r■ [l,-1] +r■ [c-1] ) ; 
120 
120<br>
121 if (L > 1.4) {<br>
122 printf ("([
 122 printf ("([Xd] : Xd,Xd,Xf) \n", rank,a, b, L); <br>123 fflush(stdout); printf (123 fflush(stdout);
\begin{array}{ccc} 124 & & & & 1 \\ 125 & & & & 1 \\ 126 & & & & 1 \\ 127 & & & & 1 \end{array}125 
1 26 
 127 
1 28 
1 29 
130 if(a>.!n) 
 131<br>
132 for (besa/berman; b++) {<br>
132 for De Ho need to calculate god if both a,b divisible by same small prime */<br>
134 ff !( ( !(aX2) & b !(bX2) ) || ( !(aX3) & b !(bX3) ) || ( !(aX5) & b !(bX5) ) || ( !(aX7) & b !(bX7) 
136 
                                           137 L•log( (doubl•)c) / (r■ [a•1]+r■Cl>·1] +r■[c-1] ); 
138 
 139 if(L > 1 .4) { 
140 printf(" ( [Xd] : ld,Xd ,Xf)\Jl", ranlt,a, b, L) ; 
141 fflu•h(atdout ); 
 142 } 
143 
 144 } 
145 } 
 146 
147 } 
148 
  149 /• Ensure all slave processors are done computing •/<br>150 MPI_Barrier(vorld);
 151 
  152 if(ramk=90) {<br>153 gettimeofday(&t_stop,&tx_dummy);<br>154 fprintf(fp,"\aTotal Time = Xg\n", diff_time(&t_start,&t_stop));<br>155 fcloss(fp);<br>printf("\n\n");<br>}
 1 58 
  159 /• IIPI c1 .. nup and ahutdovn •/ 
160 IIPI.PiD.alba( ) ; 
 161 
               return 0;
 163
```
#### 5.4.2 ut il .c



```
realloc(primes, sizeof(unsigned int));
  33 
34 
35 
36 
37 
38 
39 
40 
41 
42 
43 
44 
45 
46 
47 
48 
49 
50 
51 
52 
53 
54 
55 
56 
57 
58 
59 
60 
61 
62 
63 
64 
65 
66 
67 
68 
69 
70 
71 
72 
                           return 1; 
                    \tilde{\Sigma}priaea[0]•2; 
                    primes[1] = 3;
                     count=2;
                      i•3 ; 
p•5; 
                     while(p < n) \leqia_prim• • 1; 
ia_Wlknovn • 1; 
                               d = 3;<br>
while(is_unknovn) {<br>
a = floor (p / (float)d);<br>
if (a <= d ) {<br>
is_unknovn = 0;
                                       if ( (a•d) •• p) { 
                                               is_prime = 0;<br>is_unknovn = 0;
                                    \frac{3}{4} = d + 2;
                            \mathbf{y}if (is_prime) {<br>primes[count] = p;<br>count++;
                             F
                   \sum_{j}^{\ell}realloc( primes, count*sizeof(unsigned int));
           return count;
   73 
74 
               inline unsigned int get_r(unsigned int n, unsigned int eprimes, int num_primes) {<br>
/* Check the list of primes to see if it is good.<br>
* If k still equals 1, the number must have been<br>
* prime, therefore set k to just be th
   75 6 7 7 8 9 8 8 1 2 3 3 4 5 6 6 7 8 8 9
                       int i; 
unaign•d int It; 
                     k=1;
                       for(i=0;i<num_primes;i++) {<br>if(primes[i] > n)
                               break;<br>else if ( (n%primes[i])==0
                                    k•k•priaea[i] ; 
                     \mathfrak zif(k**1) {
   90 
91 
92 
93 
                    \frac{\text{if (k=1)}}{\text{km}}return k; 
   94 
95 
               !•--------------------------------·------.. 
/• FUNCTION diff_tilla(t1 ,t2) •/ 
/• Jleiurna a decimal value (in aeca) of elapsed ti.Jlle betv••n t1 and t2 •/ 
,------- -------------------------------, 
   96 
97 
98 
 99 
1 00 
               double diff_time(struct timeval * t_1, struct timeval * t_2) {<br>double diff;
101 
1 02 
103diff = (double) ((t_{-}2-\gt>t v_{-}sec+t_{-}2-\gt tv_{-}usec/1.0E6)- (t_1 > tv_sec+t_1->tv_usec/1.0E6));
 1 04 
1 05 
                     return (diff) ; 
106 
107 
           \overline{ }
```
# **Bibliography**

 $\sim$ 

 $\overline{\mathcal{O}}$ 

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**Further Reading: Joseph H. Silverman. The Arithmetic of Elliptic Curves., Springer-Verlag.** 

## **Vita**

Jeffrey Paul Wheeler was born April 22, 1968 in Wheeling, West Virginia. He graduated *from Linsly School in May 1986. In May 1990, Jeffrey graduated from Miami University in Oxford, Ohio earning a Bachelor of Arts degree in Mathematics with a minor in Social*  .. *Work. After graduation, Jeffrey took a position as a lecturer at Belmont Technical College in St. Clairsville, Ohio. He also served in an administrative position. In August 1998,*  Jeffrey entered the graduate program at the University of Tennessee, Knoxville. During *his four years, Jeffrey held the position of graduate teaching associate. In each of those four years, Jeffrey was a finalist for the Dorthea and Edgar Eaves Teaching Award and was recipient of the award in the academic year 2001.* 

*Upon a return to Miami Universtiy, Jeffrey met the woman who was to be his wife. The couple married June 3, 2000 in Chagrin Falls, Ohio and resided in Knoxville, Tennessee. Jeffrey's wife, Jamie, earned her MBA from the University of Tennessee, Knoxville in May 2002 while Jeffrey earned his master's degree in August 2002. Currently the couple resides in Memphis, Tennessee where Jamie is employed by FedEx and Jeffrey is pursuing a doctorate degree in Mathematics at the University of Memphis.*